

be considered as a free infinitary algebra generated by the atomic formulas and the set V of variables is supposed to be well-ordered in a sequence $x_1, x_2, \dots, x_\xi, \dots$ whose length is a cardinal $|V|$. A *language* will be identified with the set of its formulas containing finitely many free variables. A *theory* will be a set of sentences. The notion of *subformula* is defined by recursion as usual; we just notice that the subformulas of $V\langle\phi_\xi: \xi \in \mu\rangle$ are $\phi_0, \phi_1, \dots, \phi_\xi, \dots$. Following the terminology of M. Morley, a *fragment* of a language L will be a subset of L closed under: (i) formation rules of finitary first order language, including substitution of variables, and (ii) subformulas. Since the intersection of a family of fragments is again a fragment, we can talk about the fragment *generated* by a set of formulas (each containing finitely many free variables).

If Σ is a fragment of L , T is a theory included in Σ and $n \in S_\omega(V)$ (i.e., n is a finite subset of V), we let $B_{\mathbb{F}}(\Sigma, T)(n)$ be the Lindenbaum-Tarski boolean algebra of equivalent classes of formulas of Σ whose free variables belong to n . We recall that two formulas of ϕ, ψ are equivalent if $T \models (\phi \leftrightarrow \psi)$. We let $[\phi]$ be the equivalence class of ϕ .

Let Σ and Σ' be fragments of L and $T \subseteq \Sigma' \subseteq \Sigma$. A boolean homomorphism from $B_{\mathbb{F}}(\Sigma, T)(n)$ into a set algebra is a Σ' -*homomorphism* if it preserves the sums $\{[\phi_\xi]: \xi \in \mu\}$ such that $V\langle\phi_\xi: \xi \in \mu\rangle$ is a logical consequence of T belonging to Σ' . A boolean homomorphism α_n from $B_{\mathbb{F}}(\Sigma, T)(n)$ into a set algebra can be extended to an L -homomorphism of the full algebra $B_{\mathbb{F}}(L, T)(n)$ iff α_n is a Σ -homomorphism. Indeed, letting ν be the unique extension to L of the map sending an atomic formula ϕ into $\alpha_n[\phi]$, we define the extension of α_n to be the quotient of ν by T . As a corollary, any Σ -homomorphism preserves all sums of power less than μ . This corollary will be implicitly used in the sequel.

PROOF OF THE MAIN THEOREM. Let T be a theory in a $L_{\mu\omega}$ language L . Assume that every $\mu \leq \theta$. For each formula $\phi(x_1, \dots, x_k) \in L$, we associate (in a one-one fashion) a k -ary relation symbol Q_ϕ and we let L' be the $L_{\omega_1\omega}$ language generated by these symbols. We assume that L' has the same set of variables V_L of L . Let \mathcal{J}' be the unique extension to L' of the map sending $Q_\phi x_{\xi_1}, \dots, x_{\xi_k}$ into $\phi[x_{\xi_1}/x_1, \dots, x_{\xi_k}/x_k]$ for all $\phi \in L$. Furthermore, we define $\mathcal{J}: L \rightarrow L'$ by $\mathcal{J}(\phi(x_{\xi_1}, \dots, x_{\xi_k})) = Q_\phi x_{\xi_1}, \dots, x_{\xi_k}$, for all $\phi \in L$.

We define our *reduced* theory $T' = \{\sigma' \in L': T \models \mathcal{J}'(\sigma')\}$. By passage to quotients, \mathcal{J}' defines a family $\langle I'_n: n \in S_\omega(V) \rangle$ of boolean isomorphisms $I'_n: B_{\mathbb{F}}(L', T')(n) \xrightarrow{\sim} B_{\mathbb{F}}(L, T)(n)$ whose inverse, I_n , has the property that $I_n[\phi] = [\mathcal{J}(\phi)]$ for all $\phi \in L^n$ (i.e., whose free variables belong to n), by the very definition of \mathcal{J} and \mathcal{J}' .

We now show that $\mathcal{J}, \mathcal{J}'$ (or rather their restrictions to the atomic formulas) establish a one-to-one correspondence between models of T and T' with a fixed universe A .

Let \mathfrak{A}' be a model of T' with universe A . Then \mathfrak{A}' defines a family of boolean ω -homomorphisms $\alpha'_n: B_{\mathbb{F}}(L', T') \rightarrow 2^{A^n}$ by $\alpha'_n[\phi] = (\phi')^{\mathfrak{A}'} = \{a \in A^n: \mathfrak{A}' \models \phi'[a]\}$ for all $\phi' \in L'^n$. Via the (restrictions of the) canonical isomorphisms I_n , α'_n determines (by composition) a boolean ω -homomorphism $\alpha_n: B_{\mathbb{F}}(L, T)(n) \rightarrow 2^{A^n}$. By Sikorski's theorem (cf. §2), α_n is a ν -homomorphism for all $\nu \in \mu$. Using the fact that $[\mathcal{J}(\exists x_\xi \theta)] = [\exists x_\xi \mathcal{J}(\theta)]$ for all $\theta \in L$, we conclude that the structure \mathfrak{A} , obtained from \mathfrak{A}' via \mathcal{J} by interpreting the primitive n -ary relation symbol Q as